

# Generic Muchnik reducibility



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# Muchnik reducibility between structures

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## Definition

If  $\mathcal{A}$  and  $\mathcal{B}$  are countable structures, then  $\mathcal{A}$  is **Muchnik reducible** to  $\mathcal{B}$  (written  $\mathcal{A} \leq_w \mathcal{B}$ ) if every  $\omega$ -copy of  $\mathcal{B}$  computes an  $\omega$ -copy of  $\mathcal{A}$ .

- ▶  $\mathcal{A} \leq_w \mathcal{B}$  can be interpreted as saying that  $\mathcal{B}$  is intrinsically at least as complicated as  $\mathcal{A}$ .
- ▶ This is a special case of Muchnik reducibility; it might be more precise to say that the problem of presenting the structure  $\mathcal{A}$  is Muchnik reducible to the problem of presenting  $\mathcal{B}$ .
- ▶ Muchnik reducibility doesn't apply to uncountable structures.

Various approaches have been used to extend computable structure theory beyond the countable:

- ▶ Computability on admissible ordinals (aka  $\alpha$ -recursion theory)
- ▶ Computability on separable structures, as in computable analysis
- ▶ ...

## Generic Muchnik reducibility

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Noah Schweber extended Muchnik reducibility to arbitrary structures (see Knight, Montalbán, Schweber):

### Definition (Schweber)

If  $\mathcal{A}$  and  $\mathcal{B}$  are (possibly uncountable) structures, then  $\mathcal{A}$  is **generically Muchnik reducible** to  $\mathcal{B}$  (written  $\mathcal{A} \leq_w^* \mathcal{B}$ ) if  $\mathcal{A} \leq_w \mathcal{B}$  in some forcing extension of the universe in which  $\mathcal{A}$  and  $\mathcal{B}$  are countable.

It follows from Shoenfield absoluteness that generic Muchnik reducibility is robust.

### Lemma (Schweber)

If  $\mathcal{A} \leq_w^* \mathcal{B}$ , then  $\mathcal{A} \leq_w \mathcal{B}$  in *every* forcing extension that makes  $\mathcal{A}$  and  $\mathcal{B}$  countable.

Note that for countable structures,  $\mathcal{A} \leq_w^* \mathcal{B} \iff \mathcal{A} \leq_w \mathcal{B}$ .

## Initial example

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### Definition (Cantor space)

Let  $\mathcal{C}$  be the structure with universe  $2^\omega$  and predicates  $P_n(X)$  that hold if and only if  $X(n) = 1$ .

### Observation (Knight, Montalbán, Schweber)

$$\mathcal{C} \leq_w^* (\mathbb{R}, +, \cdot).$$

To understand this example, say that we take a forcing extension that collapses the continuum.

The Turing degrees from the ground model now form a countable ideal  $I$ . By absoluteness, this ideal has many of the properties it has in the ground model. It's a jump ideal and much more.

Let  $\mathbb{R}_I$  be the reals in  $I$  (the ground model's version of  $\mathbb{R}$ ). Similarly, let  $\mathcal{C}_I$  denote the restriction of  $\mathcal{C}$  to sets in  $I$  (the ground model's version of  $\mathcal{C}$ ).

## Initial example

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### Facts

- ▶ From a copy of  $(\mathbb{R}_I, +, \cdot)$ , or even  $(\mathbb{R}_I, +, <)$ , we can compute an *injective* listing of the sets in  $I$ , i.e., one with no repetitions.
- ▶ A degree  $\mathbf{d}$  computes a copy of  $\mathcal{C}_I$  iff it computes an (injective) listing of the sets in  $I$ .

This shows that  $\mathcal{C}_I \leq_w (\mathbb{R}_I, +, <)$ . It is even easier to see that  $(\mathbb{R}_I, +, <) \leq_w (\mathbb{R}_I, +, \cdot)$ .

Therefore,  $\mathcal{C} \leq_w^* (\mathbb{R}, +, <) \leq_w^* (\mathbb{R}, +, \cdot)$ .

### Question (Knight, Montalbán, Schweber)

$$\text{Is } (\mathbb{R}, +, \cdot) \leq_w^* \mathcal{C}?$$

No! This was answered by Igusa and Knight, and independently (though later) by Downey, Greenberg, and M.

## Facts about $\mathcal{C}$ and $\mathcal{B}$

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### Definition (Baire space)

Let  $\mathcal{B}$  be the structure with universe  $\omega^\omega$  and, for each finite string  $\sigma \in \omega^{<\omega}$ , a predicate  $P_\sigma(f)$  that holds if and only if  $\sigma < f$ .

The following facts were proved by Igusa, Knight; Downey, Greenberg, M.; Igusa, Knight, Schweber; Andrews, Knight, Kuyper, Lempp, M., Soskova.

- ▶  $\mathcal{B} \equiv_w^* (\mathbb{R}, +, <) \equiv_w^* (\mathbb{R}, +, \cdot)$ . This degree also contains every closed/continuous expansion of  $(\mathbb{R}, +, \cdot)$ .
- ▶  $\mathcal{C} <_w^* \mathcal{B}$ .
- ▶  $\mathcal{C}' \equiv_w^* \mathcal{B}$ .
- ▶ The closed expansions of  $\mathcal{C}$  lie in the interval between  $\mathcal{C}$  and  $\mathcal{B}$ .

### Question

Is there a generic Muchnik degree strictly between  $\mathcal{C}$  and  $\mathcal{B}$ ?

## Definability and post-extension complexity

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It is going to be important to understand the complexity of definable sets both before and after the forcing extension.

### Definition

We say that a relation  $R$  on a structure  $\mathcal{M}$  is  $\Sigma_n^c(\mathcal{M})$  if it is definable by a computable  $\Sigma_n$  formula in  $\mathcal{L}_{\omega_1\omega}$  with finitely many parameters.

### Theorem (Ash, Knight, Manasse, Slaman; Chisholm)

If  $\mathcal{M}$  is countable, then  $R$  is  $\Sigma_n^c(\mathcal{M})$  if and only if it is relatively intrinsically  $\Sigma_n^0$ , i.e., its image in any  $\omega$ -copy of  $\mathcal{M}$  is  $\Sigma_n^0$  relative to that copy.

Computable objects and satisfaction on a structure are absolute, so:

### Corollary

A relation  $R$  is  $\Sigma_n^c(\mathcal{M})$  if and only if it is relatively intrinsically  $\Sigma_n^0$  in any/every forcing extension that makes  $\mathcal{M}$  countable.

## Definability and pre-extension complexity

In structures like  $\mathcal{C}$  and  $\mathcal{B}$ , we can also measure the complexity of  $\Sigma_n^c(\mathcal{M})$  relations in topological terms.

The calculation depends on the structure:

	$\Sigma_2^c$	$\Sigma_3^c$	$\Sigma_4^c$	$\Sigma_5^c$	$\Sigma_6^c$	...
$\mathcal{B}$	$\Sigma_1^1$	$\Sigma_2^1$	$\Sigma_3^1$	$\Sigma_4^1$	$\Sigma_5^1$	...
$\mathcal{C}$	$\Sigma_2^0$	$\Sigma_1^1$	$\Sigma_2^1$	$\Sigma_3^1$	$\Sigma_4^1$	...

- ▶ These bounds are sharp, e.g., every  $\Sigma_1^1$  relation on  $\mathcal{B}$  is  $\Sigma_2^c(\mathcal{B})$ .
- ▶ The “lost quantifiers” correspond to the first order quantifiers needed in the normal form for  $\Sigma_n^1$  relations with function/set quantifiers.
- ▶ This leads to an easy (and essentially different) separation between the generic Muchnik degrees of  $\mathcal{C}$  and  $\mathcal{B}$ .



## A degree strictly between $\mathcal{C}$ and $\mathcal{B}$ (ver. 1.0)

### Lemma

There is a linear order  $\mathcal{L}$  such that  $\mathcal{L} \leq_w^* \mathcal{B}$  but  $\mathcal{L} \not\leq_w^* \mathcal{C}$ .

Idea: code a  $\Pi_2^1$  complete set into  $\mathcal{L}$  so that it can be extracted in a  $\Sigma_4^c$  way.

### Lemma

If  $\mathcal{L}$  is a linear order, then  $\mathcal{B} \not\leq_w^* \mathcal{C} \sqcup \mathcal{L}$ .

Similar to the Downey, Greenberg, M. proof that  $\mathcal{B} \not\leq_w^* \mathcal{C}$ ; we show that a generic countable presentation of  $\mathcal{C} \sqcup \mathcal{L}$  does not compute a copy of  $\mathcal{B}$ . The key fact used about linear orders is that their  $\sim_2$ -equivalence classes are tame (Knight 1986).

Now let  $\mathcal{M} = \mathcal{C} \sqcup \mathcal{L}$ , where  $\mathcal{L}$  is the linear order from the first lemma.

### Corollary

There is a structure  $\mathcal{M}$  such that  $\mathcal{C} <_w^* \mathcal{M} <_w^* \mathcal{B}$ .

## Degrees strictly between $\mathcal{C}$ and $\mathcal{B}$ (ver. 2.0)

Joining  $\mathcal{C}$  with the right linear order was a (somewhat awkward) way of making a new set  $\Sigma_4^c$  definable (without lifting us up to  $\mathcal{B}$ ).

There is a more natural way to do this:

### Theorem (Gura)

Using marker extensions, we can build structures

$$\mathcal{C} <_w^* \dots <_w^* \mathcal{M}_3 <_w^* \mathcal{M}_2 <_w^* \mathcal{M}_1 <_w^* \mathcal{B}$$

with the following “complexity profiles”:

	$\Sigma_2^c$	$\Sigma_3^c$	$\Sigma_4^c$	$\Sigma_5^c$	$\Sigma_6^c$	...
$\mathcal{B}$	$\Sigma_1^1$	$\Sigma_2^1$	$\Sigma_3^1$	$\Sigma_4^1$	$\Sigma_5^1$	...
$\mathcal{M}_1$	$\Sigma_2^0$	$\Sigma_2^1$	$\Sigma_3^1$	$\Sigma_4^1$	$\Sigma_5^1$	...
$\mathcal{M}_2$	$\Sigma_2^0$	$\Sigma_1^1$	$\Sigma_3^1$	$\Sigma_4^1$	$\Sigma_5^1$	...
$\mathcal{M}_3$	$\Sigma_2^0$	$\Sigma_1^1$	$\Sigma_2^1$	$\Sigma_4^1$	$\Sigma_5^1$	...
			$\vdots$			
$\mathcal{C}$	$\Sigma_2^0$	$\Sigma_1^1$	$\Sigma_2^1$	$\Sigma_3^1$	$\Sigma_4^1$	...

## Open questions

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1. Can an *expansion* of  $\mathcal{C}$  be strictly between  $\mathcal{C}$  and  $\mathcal{B}$ ?
2. Are the degrees of  $\mathcal{M}_1, \mathcal{M}_2, \mathcal{M}_3, \dots$  the only degrees strictly between  $\mathcal{C}$  and  $\mathcal{B}$ ?
3. Are there incomparable degrees between  $\mathcal{C}$  and  $\mathcal{B}$ ?

## Expansions of $\mathcal{C}$ above $\mathcal{B}$

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Let  $\mathcal{M} = (\mathcal{C}, \text{Stuff})$  be an expansion of  $\mathcal{C}$ . First, we want a criterion that guarantees that  $\mathcal{M} \geq_w^* \mathcal{B}$ .

- ▶ If the set  $\mathcal{F} \subset 2^\omega$  of sequences with finitely many ones is  $\Delta_1^c(\mathcal{M})$ , i.e., computable in every  $\omega$ -copy of  $\mathcal{M}$ , then  $\mathcal{M} \geq_w^* \mathcal{B}$ .
  - ▶ Why? There is a natural bijection between  $\mathcal{B}$  and  $\mathcal{C} \setminus \mathcal{F}$ .
- ▶ If  $\mathcal{F}$  is  $\Delta_2^c(\mathcal{M})$ , then  $\mathcal{M} \geq_w^* \mathcal{B}$ .
  - ▶ Add a little injury.
  - ▶ This lets us show, for example, that  $(\mathcal{C}, \oplus) \geq_w^* \mathcal{B}$ .
- ▶ If any countable dense set is  $\Delta_2^c(\mathcal{M})$ , then  $\mathcal{M} \geq_w^* \mathcal{B}$ .
- ▶ If there is a perfect set  $\mathcal{P} \subseteq \mathcal{C}$  with a countable dense  $\mathcal{Q} \subset \mathcal{P}$  that is  $\Delta_2^c(\mathcal{M})$ , then  $\mathcal{M} \geq_w^* \mathcal{B}$ .

## Expansions of $\mathcal{C}$ above $\mathcal{B}$

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- ▶ If there is a perfect set  $\mathcal{P} \subseteq \mathcal{C}$  with a countable dense  $Q \subset \mathcal{P}$  that is  $\Delta_2^c(\mathcal{M})$ , then  $\mathcal{M} \geq_w^* \mathcal{B}$ .

### Lemma

If  $\mathcal{M} \leq_w^* \mathcal{B}$  and  $R \subseteq \mathcal{C}$  is  $\Delta_2^c(\mathcal{M})$ , then it is  $\Delta_2^c(\mathcal{B})$ , i.e., Borel.

### Lemma (Hurewicz)

If  $R \subseteq \mathcal{C}$  is Borel but not  $\Delta_2^0$ , then there is a perfect set  $\mathcal{P} \subseteq \mathcal{C}$  such that either  $\mathcal{P} \cap R$  or  $\mathcal{P} \setminus R$  is countable and dense in  $\mathcal{P}$ .

Putting it all together (and noting that arity doesn't matter):

### Lemma

If  $\mathcal{M} \leq_w^* \mathcal{B}$  is an expansion of  $\mathcal{C}$  and  $R \subseteq \mathcal{C}^n$  is  $\Delta_2^c(\mathcal{M})$  but not  $\Delta_2^0$ , then  $\mathcal{M} \geq_w^* \mathcal{B}$ .

# Tameness and dichotomy

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In the contrapositive (and using the fact that  $\Delta_2^0 = \Delta_2^c(\mathcal{C})$ ):

## Tameness Lemma

If  $\mathcal{M} <_w^* \mathcal{B}$  is an expansion of  $\mathcal{C}$ , then  $\Delta_2^c(\mathcal{M}) = \Delta_2^c(\mathcal{C})$ .

## Dichotomy Theorem for Closed Expansions

If  $\mathcal{M} \leq_w^* \mathcal{B}$  is an expansion of  $\mathcal{C}$  by closed relations (and/or continuous functions), then either  $\mathcal{M} \equiv_w^* \mathcal{C}$  or  $\mathcal{M} \equiv_w^* \mathcal{B}$ .

Combined with work of Greenberg, Igusa, Turetsky, and Westrick:

## Dichotomy Theorem for Unary Expansions

If  $\mathcal{M} \leq_w^* \mathcal{B}$  is an expansion of  $\mathcal{C}$  by countably many unary relations, then either  $\mathcal{M} \equiv_w^* \mathcal{C}$  or  $\mathcal{M} \equiv_w^* \mathcal{B}$ .

These dichotomy results take care of most natural (and many unnatural) examples of expansions.

## Open questions

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1. Can an *expansion* of  $\mathcal{C}$  be strictly between  $\mathcal{C}$  and  $\mathcal{B}$ ? (In particular, the non-unary  $\Delta_2^0$  case is open.)
2. Are the degrees of  $\mathcal{M}_1, \mathcal{M}_2, \mathcal{M}_3, \dots$  the only degrees strictly between  $\mathcal{C}$  and  $\mathcal{B}$ ?
3. Are there incomparable degrees between  $\mathcal{C}$  and  $\mathcal{B}$ ?

These questions are related. For example:

**Fact.** Any Borel expansion of  $\mathcal{C}$  that is not above  $\mathcal{B}$  has the same complexity profile as  $\mathcal{C}$ . So a positive answer to 1 gives a negative answer to 2.

We have focused on  $\mathcal{C}$  and  $\mathcal{B}$  (and a couple of other degrees). What else are generic Muchnik degrees good for?

THANK YOU!