

CS/ECE/MATH 435, HOMEWORK 10, DUE MAY 8.

1. Which of the following is not traditionally an information source for authenticating someone's identity? Explain.

- (a) Something you know.
- (b) Something you have.
- (c) Something you like.
- (d) Something you are.

2. You are given a piece of data. Storage is limited so you also need to compress the data. You are given RSA for encryption and signing and a good compression algorithm (e.g. LZW).

- (a) Should you encrypt first or compress? Or does the order not matter? Why?
- (b) Should you sign first or compress? Or does the order not matter? Why?

3. Zero-knowledge proofs are where you convince someone you can do something without actually giving away the proof. For example, suppose Alice wants to convince Bob that she knows a number x without Bob figuring out x (this has applications e.g. in banking).

Here's how she does it. She picks two large primes p, q and sets $N = pq$. She picks a number x between 1 and N . She tells Bob N and $x^2 \pmod{N}$ (over a public channel). If Bob could factor N , then he could compute x and it is believed that there is no easier way to find x .

Alice now picks a random integer r and sends Bob $x^2 r^2 \pmod{N}$. Bob randomly sends one of two questions - "Send me r " or "Send me $xr \pmod{N}$ ".

- (a) Show that Alice can satisfy both these requests.
- (b) Show that Bob can check her answer in either case.
- (c) Suppose Oscar tries to fool Bob by making up a random number s and sending s^2 to Bob. Show that if Bob asks for $xr \pmod{N}$, Oscar is OK, but that if Bob asks for r , then Oscar is caught. Why does this mean that by playing this game several times with different r , Alice gives a zero-knowledge proof with high probability.

[There is an algorithm that given r and $xr \pmod{N}$ lets you calculate x .]