

Division
Algorithm

quotient &
remainder

(repeated
division
algorithm)

↓ lemma

Euclid's
Algorithm

gcd

- greatest common
divisor

Bezout's
Identity

gcd as a
linear combination

(work backward)
(find an inverse
of $a \text{ mod } m$)

$ax \equiv b \pmod{m}$

Solving linear
congruence relations

Chinese Remainder
Theorem

Solving SYSTEMS
of linear
congruence
relations

Application

Computer Arithmetic
with Large Numbers

A Note on Notation: Typically we prefer the notation "1" for "such that" in sets $\{ \quad | \quad \}$ however, we do not wish to confuse this vertical line for an instance of "divides" which we also write using a vertical line. Therefore if either $\boxed{\quad}$ or \blacksquare contain a "divides" symbol we use a colon ":" to express the phrase "such that."

Greatest Common Divisors

let $a, b \in \mathbb{N}$

$$\text{Div}(a) := \{x : x \mid a\} \quad (\text{set of } \underline{\text{divisors}} \text{ of } a)$$

$$\text{Div}(b) := \{x : x \mid b\} \quad (\text{set of } \underline{\text{divisors}} \text{ of } b)$$

$$1 \in \overbrace{\text{Div}(a) \cap \text{Div}(b)}^{\neq \emptyset \text{ non-empty}} \quad (\text{set of } \underline{\text{common divisors}} \text{ of } a \& b)$$

$$\gcd(a, b) := \max(\text{Div}(a) \cap \text{Div}(b))$$

(the greatest common divisor)

e.g. $\text{Div}(12) = \{1, 2, 3, 4, 6, 12\}$

$$\text{Div}(18) = \{1, 2, 3, 6, 9, 18\}$$

$$\text{Div}(12) \cap \text{Div}(18) = \{1, 2, 3, 6\}$$

$$\max(\text{Div}(12) \cap \text{Div}(18)) = 6$$

$$\Rightarrow \gcd(12, 18) = 6$$

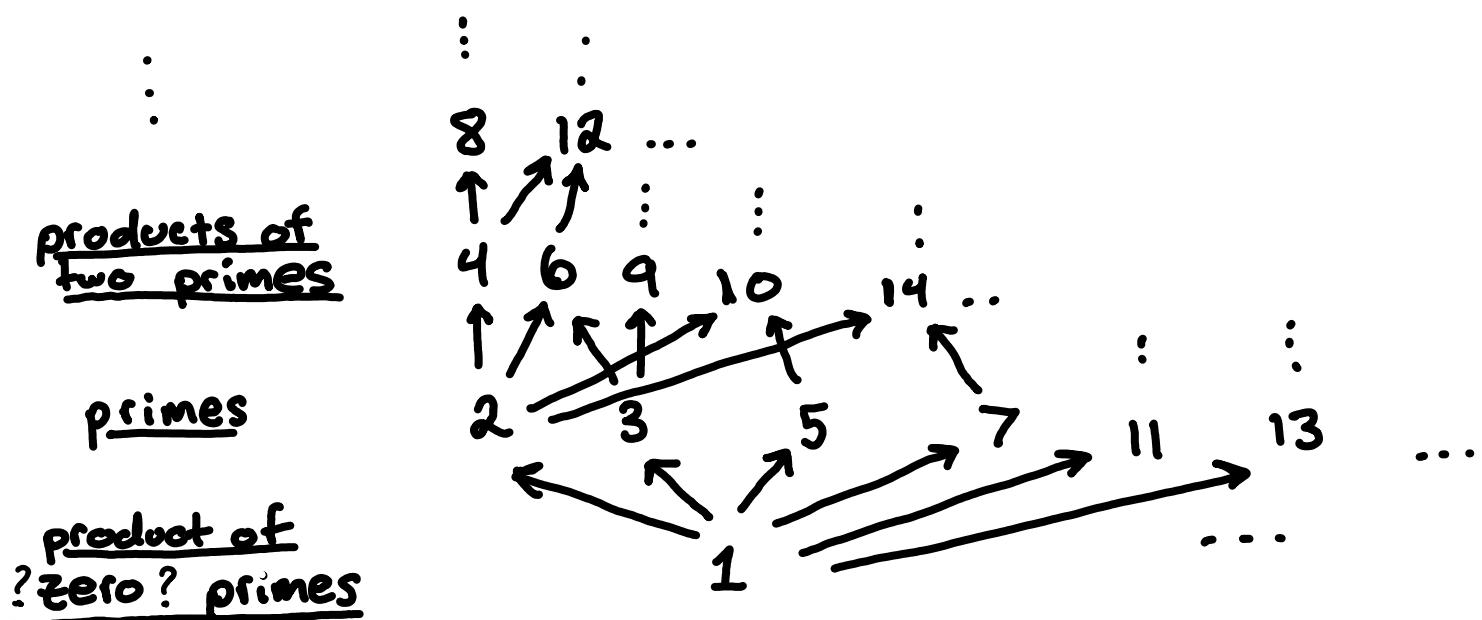
Note: If $X \subseteq \mathbb{N}$, and $X \neq \emptyset$, $|X| < \infty$ then $\max X$ exists.

A Second approach to GCDs

Recall that "a divides b" defines a relation on the natural numbers \mathbb{N} . What properties did it have?

- pre-order [(i) Reflexive: $a|a$ since $a = 1 \cdot a$]
partial order [(ii) Transitive: $a|b \wedge b|c \Rightarrow a|c$
since $ax = b \wedge by = c$
implies $a(xy) = c$]
order [(iii) Anti-Symmetric: $a|b \wedge b|a \Rightarrow a = b$
 \Downarrow
 $a \leq b \wedge b \leq a \Rightarrow$]

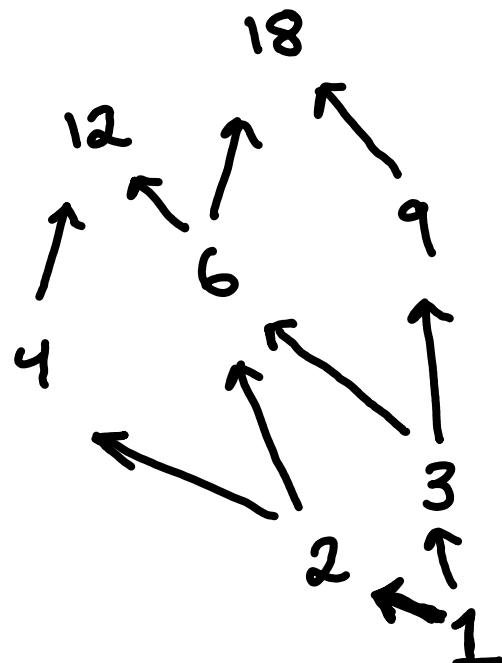
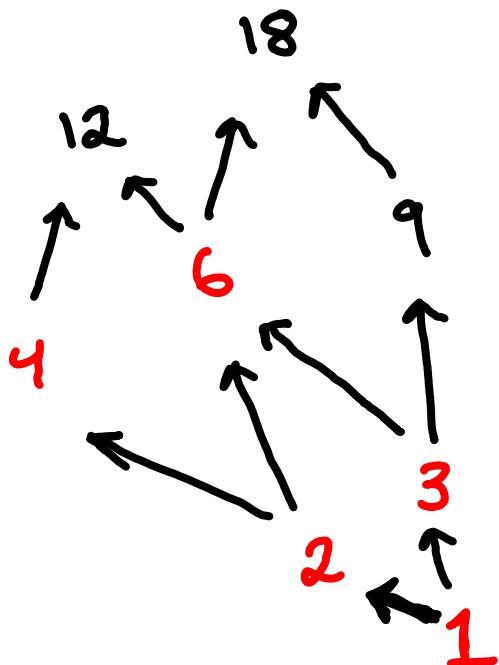
Thus, we can draw the Hasse diagram of the poset $(\mathbb{N}, |)$ & hope to understand what GCDs are geometrically (i.e. pictorially). Here is a small portion of the diagram.



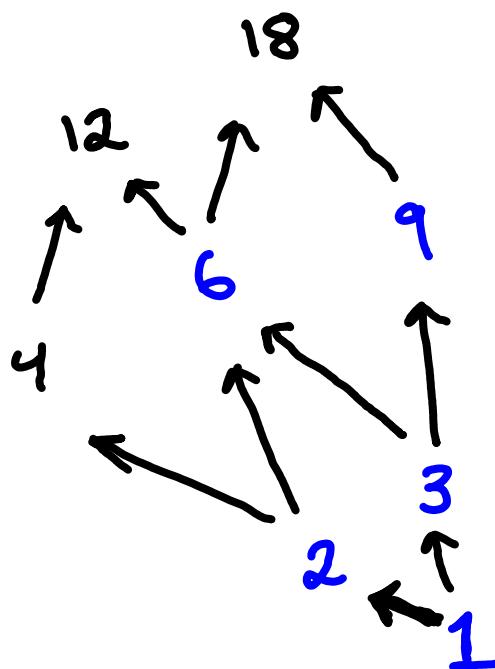
this picture gets crowded and unwieldy rather quickly, so it is useful to omit parts deemed not immediately relevant to the task at hand.

Here is the minimal piece of the diagram needed to understand the calculation of $\gcd(12, 18)$.

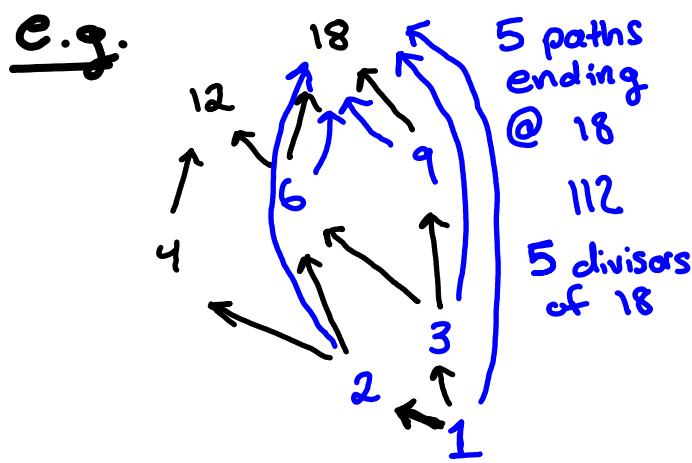
Coloring all divisors of 12 **red** we see



Similarly, we might color the divisors of 18 **blue** to get

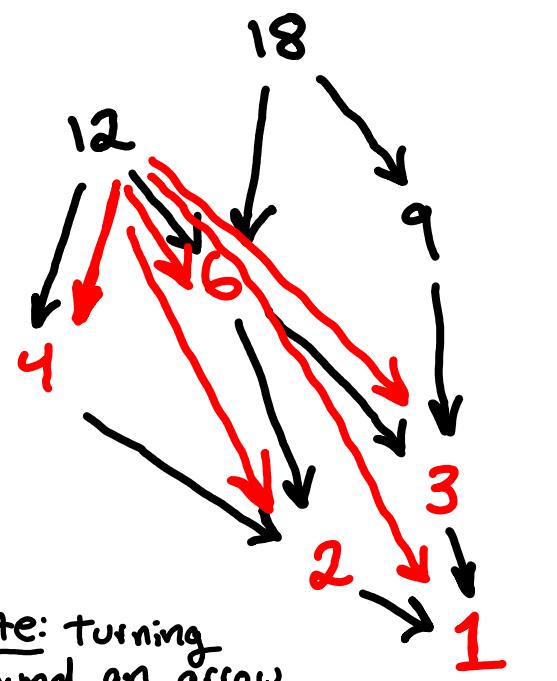


It is not difficult to "see" the set of divisors of a number in this poset, indeed, $\text{Div}(n) = \{\text{all numbers that can reach } n \text{ travelling along arrows of the diagram}\}$. To say this in another way, $\text{Div}(n)$ is the collection of numbers appearing as the starting point of a path (of directed arrows) ending @ n .

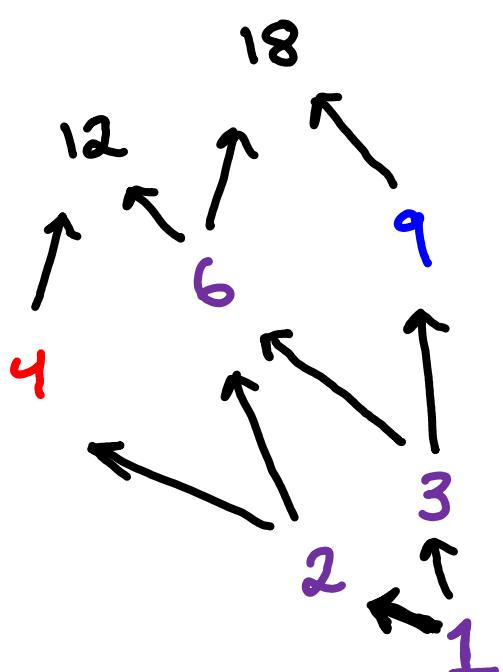


To shorten what must be said, we can summarize as follows
 $\text{Div}(n) = \{\text{numbers pointing to } n\}$.
 One might also choose to turn the arrows around, in which case
 $\text{Div}(n) = \{n \text{ points to}\}$

If we color the divisors of 12 red, and at the same time color the divisors of 18 blue, letting two colors mix to form purple when they coincide, we get.



Note: turning around an arrow indicating a "divides" relation results in an arrow indicating the relation "is a multiple of"



$$\text{Div}(12) \cap \text{Div}(18)$$

By definition, the purple numbers are common divisors of 12 & 18.

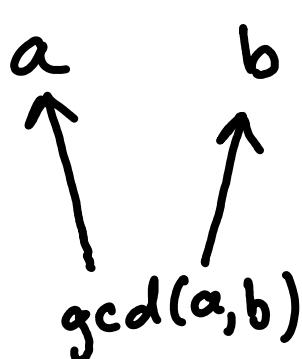
$$\begin{aligned} \text{Div}(12) &= \{x : x \mid 12\} \\ &\quad // \\ &= \{x \mid x \rightarrow 12\} // \end{aligned}$$

$$\begin{aligned} \text{Div}(18) &= \{x : x \mid 18\} \\ &\quad // \\ &= \{x \mid x \rightarrow 18\} // \end{aligned}$$

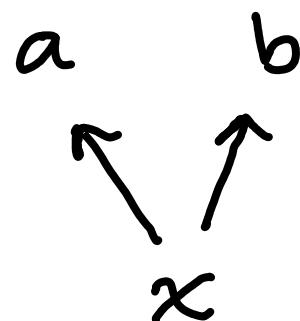
Lastly, we must figure out how to "see" which of the purple numbers is largest from the Hasse diagram. By positioning numbers higher on the page when they are larger in absolute value, we can easily spot this maximum of the numbers colored purple earlier i.e., the GCD

Since the purple number farthest up the page is the GCD (of the numbers in question, 12 & 18 in our case) & any purple number is a divisor of both 12 & 18, all purple numbers point to the one furthest up the page!

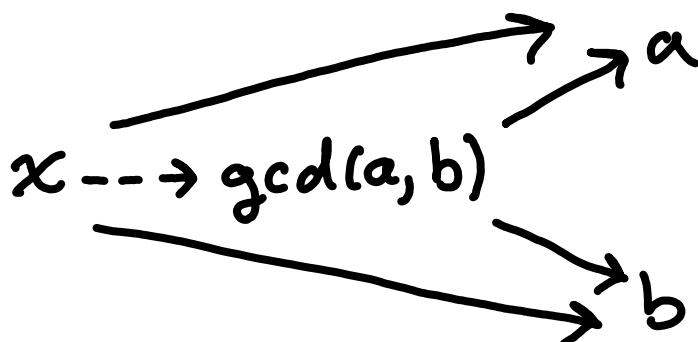
$\gcd(a, b)$ is said to be "the last number pointing to both a & b " in the sense that



& whenever



one has $x \rightarrow \gcd(a, b)$. Concisely,



the dashed arrow is the one guaranteed by "context" (the context being the solid arrows). The origin of this description of GCDs is category theory. If you are curious, the relevant terms to look up are "product" or "limit." Turning all arrows around produces the definition of a "coproduct" or "colimit."

Q: What happens if we turn the arrows around in the poset $(\mathbb{N}, |)$ then find the last number

pointing to a & b ? What is the relationship between this number & $\gcd(a, b)$?

These questions will be answered at another time. For now, we pave the way for this future discussion with an observation

Observe: Organizing the naturals by height on the page according to the # of prime factors (counting with multiplicity so $4 = 2^2$ has two prime factors) the last row with at least one purple number has at most one purple number.

proof sketch: If two numbers x & y have the same number of prime factors and $x \mid y$ then the prime factorization of x coincides with that of y and thus $x = y$. The observation follows from setting $y = \gcd(a, b)$ & x any common divisor of a & b at the same height as y in the poset (\mathbb{N}, \mid) described above. \square

It is then reasonable to guess

$$\gcd(a, b) = \begin{matrix} \text{product of all primes} \\ a \& b \text{ have in common} \\ \text{with multiplicity.} \end{matrix}$$

Indeed this is what we will eventually prove.

The following are equivalent

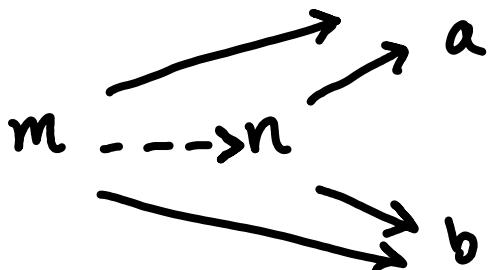
(i) n is the greatest common divisor of a & b i.e., $\gcd(a, b) = n$

(ii) $\max(\text{Div}(a) \cap \text{Div}(b)) = \{n\}$.

(iii) n is the largest number to divide both a & b i.e., $n|a$ & $n|b$ and $\forall m |a$ & $m|b \Rightarrow m|n$.

(iv) n is the last thing pointing into both a & b in the poset (\mathbb{N}, \mid) i.e., $m \rightarrow a$ & $m \rightarrow b \Rightarrow m \rightarrow n$

in category theory the GCD of a & b is the product of a & b in the poset category (\mathbb{N}, \mid)



Such an n is called the "infimum" or "meet" of a & b in the poset (\mathbb{N}, \mid)

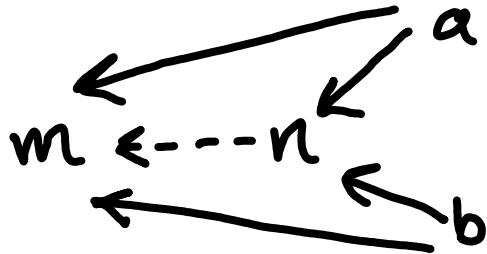
(v) n is the product of all primes common to both a & b counted with multiplicity

$$\text{i.e., } a = p_1^{a_1} p_2^{a_2} \dots p_n^{a_n}, b = p_1^{b_1} p_2^{b_2} \dots p_n^{b_n}$$

$$\text{then } n = p_1^{\min(a_1, b_1)} p_2^{\min(a_2, b_2)} \dots p_n^{\min(a_n, b_n)}$$

(vii) n is the first thing both a & b point into if the relation defining the poset (\mathbb{N}, \mid) is reversed i.e., if $x \rightarrow y$ if and only if $y \mid x$ then

In Category theory
(if you are curious)
"turning around arrows"
is called "dualization" &
the result is the "opposite
category" & the diagram
to the right defines n
as the "coproduct" of
 a & b or the "colimit"
of the diagram consisting
of a & b .



Such an n is called the
"supremum" or "join" of a & b
in the poset (\mathbb{N}, \mid) with the
relation reversed, i.e., the
inverse relation of \mid on \mathbb{N} .

lastly, we restate the unanswered question from earlier before moving on.

Q: What would happen if we were to find
the join of a & b in (\mathbb{N}, \mid) without first
taking the inverse relation (or equivalently, reversing
the direction of arrows in the Hasse diagram)?

Road Map

prerequisite

GCDs



Next:

Lemma

Division
Algorithm

Euclidean
Algorithm

Lemma: Suppose $a = qb + r$ w/ $0 \leq r < b$

then $\gcd(a, b) = \gcd(b, r)$

proof: Show that the set of common divisors agree in both cases

Suppose $d | a$ & $d | b$ } $\text{Div}(a) \cap \text{Div}(b)$
then $d | r = a - qb$ } \cap (is a subset of)
 $\text{Div}(b) \cap \text{Div}(r)$

Conversely, if $d | b$ & $d | r$ } $\text{Div}(b) \cap \text{Div}(r)$
then $d | a = qb + r$ } \cap
 $\text{Div}(a) \cap \text{Div}(b)$

Hence $\text{Div}(a) \cap \text{Div}(b) = \text{Div}(b) \cap \text{Div}(r)$

\Rightarrow they have the same maximum element

$\therefore \gcd(a, b) = \gcd(b, r)$ \square

Euclid's Algorithm

Q: What is $\gcd(4620, 101)$?

1) Division Algorithm $\rightarrow q_1$ & r_1

2) Use lemma to find gcd of smaller #s

3) Division Algorithm $\rightarrow q_2$ & r_2

:

$$4620 = 45 \cdot 101 + 75$$

$$101 = 1 \cdot 75 + 26$$

$$75 = 2 \cdot 26 + 23$$

$$26 = 1 \cdot 23 + 3$$

$$23 = 7 \cdot 3 + 2$$

$$3 = 1 \cdot 2 + 1$$

$$2 = 2 \cdot 1 + 0$$

← Stop once there is no remainder.

gcd is remainder from second to last line

$$\begin{aligned} A: \quad \text{gcd}(4620, 101) &= \text{gcd}(101, 75) \\ &= \text{gcd}(75, 26) = \text{gcd}(26, 23) \\ &= \text{gcd}(23, 3) = \text{gcd}(3, 2) = 1 \end{aligned}$$

Def: a and b are said to be relatively prime if $\text{gcd}(a, b) = 1$

Hence 4620 & 101 are relatively prime.

Bezout's Identity

$a, b \in \mathbb{Z}_{>0} \Rightarrow \exists s, t \in \mathbb{Z}$ s.t

$$\gcd(a, b) = sa + tb$$

$$\begin{aligned}
 4620 &= 45 \cdot 101 + 75 \\
 101 &= 1 \cdot 75 + 26 \\
 75 &= 2 \cdot 26 + 23 \\
 26 &= 1 \cdot 23 + 3 \\
 23 &= 7 \cdot 3 + 2 \\
 3 &= 1 \cdot 2 + 1 \\
 2 &= 2 \cdot 1 + 0
 \end{aligned}$$

flip
→

$$\begin{aligned}
 2 &= 2 \cdot 1 + 0 && \text{forget first line} \\
 3 &= 1 \cdot 2 + 1 \\
 23 &= 7 \cdot 3 + 2 \\
 26 &= 1 \cdot 23 + 3 \\
 75 &= 2 \cdot 26 + 23 \\
 101 &= 1 \cdot 75 + 26
 \end{aligned}$$

$$4620 = 45 \cdot 101 + 75$$

$$3 = 1 \cdot 2 + 1 \xrightarrow{\text{gcd rearrange}} 1 = 3 - 1 \cdot 2$$

$$23 = 7 \cdot 3 + 2 \xrightarrow{} 2 = 23 - 7 \cdot 3$$

$$26 = 1 \cdot 23 + 3 \xrightarrow{} 3 = 26 - 1 \cdot 23$$

$$75 = 2 \cdot 26 + 23 \xrightarrow{} 23 = 75 - 2 \cdot 26$$

$$101 = 1 \cdot 75 + 26 \xrightarrow{} 26 = 101 - 1 \cdot 75$$

$$4620 = 45 \cdot 101 + 75 \xrightarrow{} 75 = 4620 - 45 \cdot 101$$

$$\begin{aligned}
 1 &= 3 - 1 \cdot 2 \\
 2 &= 23 - 7 \cdot 3 \\
 3 &= 26 - 1 \cdot 23 \\
 23 &= 75 - 2 \cdot 26 \\
 26 &= 101 - 1 \cdot 75 \\
 75 &= 4620 - 45 \cdot 101
 \end{aligned}$$

$$\begin{aligned}
 1 &= 3 - 1 \cdot (23 - 7 \cdot 3) \\
 &= 8 \cdot 3 - 1 \cdot 23 \\
 &= 8 \cdot (26 - 1 \cdot 23) - 1 \cdot 23 \\
 &= 8 \cdot 26 - 9 \cdot 23 \\
 &= 8 \cdot 26 - 9 \cdot (75 - 2 \cdot 26) \\
 &= 26 \cdot 26 - 9 \cdot 75 \\
 &= 26 \cdot (101 - 1 \cdot 75) - 9 \cdot 75 \\
 &= 26 \cdot 101 - 35 \cdot 75 \\
 &= 26 \cdot 101 - 35 \cdot (4620 - 45 \cdot 101) \\
 1 &= 1601 \cdot 101 - 35 \cdot 4620
 \end{aligned}$$

Express remainders as linear combinations
& replace this expression in the line above it

Linear Congruence

Solve for x if

$$ax \equiv b \pmod{m}$$

Note: There may not be a solution

e.g.

$$7x \equiv 1 \pmod{7}$$

has no solution since multiples of 7
are congruent to 0 mod 7.

Q: When Can we guarantee a solution?

A: If $\gcd(a, m) = 1$ (i.e. a & m are relatively prime) Then $\exists!$ Solution to the linear congruence

$$ax \equiv b \pmod{m}$$

The main idea:

How would we find the answer to $ax = b$?
we hope $a \neq 0$ so we may divide by it
to get $a^{-1}b = (a^{-1}a)x = 1 \cdot x = x$.

Solve a simpler version of the problem then try to modify the solution to work in the original situation.

Multiplicative Inverses Mod m

If $\bar{a}a \equiv 1 \pmod{m}$

then we call \bar{a} an inverse of a modulo m . This is the analogue of dividing by a in modular arithmetic.

Theorem: $\gcd(a, m) = 1 \Rightarrow \exists!$ inverse of a modulo m .

Proof: $\gcd(a, m) = 1 \Rightarrow \exists s, t$ such that

$sa + tm = 1$ Reduce both sides
of the equation modub m to see

$$sa \equiv 1 \pmod{m} \text{ Hence}$$

s is the inverse of a modulo m . \square

The Chinese Remainder Theorem

e.g. Solve for x if

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

Spoiler Alert: $x = 23$ works because

$$23 \bmod 3 = 2 \quad (\text{remainder after division})$$

$$23 \bmod 5 = 3$$

$$\& \quad 23 \bmod 7 = 2$$

Method 1 (back substitution) \leftarrow Not the CRT

$$x \equiv 2 \pmod{3} \iff \exists k \in \mathbb{Z} \text{ s.t. } x = 3k + 2$$

$$x \equiv 3 \pmod{5} \iff 3k+2 \equiv 3 \pmod{5}$$

(Subtract 2 from both sides) $3k \equiv 1 \pmod{5}$

($\gcd(3,5)=1 \Rightarrow$ can solve for k) $k \equiv 2 \pmod{5}$

$$2 \cdot 3 - 5 = 1$$



$$\exists \tilde{k} \in \mathbb{Z} \text{ s.t. } k = 5\tilde{k} + 2$$

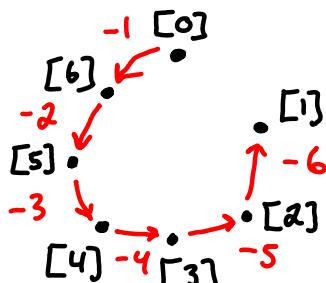
$$\begin{aligned} \Rightarrow x &= 3k + 2 = 3(5\tilde{k} + 2) + 2 \\ &= 15\tilde{k} + 6 + 2 = \underline{\underline{15\tilde{k} + 8}} \end{aligned}$$

$x \equiv 2 \pmod{7}$ Finally, we use the third congruence relation

$$15\tilde{k} + 8 \equiv 2 \pmod{7} \quad \text{then solve}$$

$$15\tilde{k} \equiv -6 \pmod{7}$$

$$15\tilde{k} \equiv 1 \pmod{7}$$



$$\gcd(15, 7) = 1 \quad \& \quad 1 = 15 - 2 \cdot 7$$

Thus $\tilde{x} \equiv 1 \pmod{7} \Leftrightarrow \tilde{x} = 7\hat{k} + 1$

for some integer \hat{k} . Thus

$$x = 15\tilde{x} + 8 = 15(7\hat{k} + 1) + 8 = 105\hat{k} + 23$$

We now have an expression for All solutions to the original system of linear congruences.

$$x \equiv 23 \pmod{105}$$

\Rightarrow possible answers are

$$\begin{array}{c} 23 \leftarrow \text{Smallest answer} \\ 128 \\ 233 \\ \vdots \end{array}$$

Note:

$$105 = 3 \cdot 5 \cdot 7$$

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

Theorem: (CRT)

If m_1, m_2, \dots, m_n are pairwise relatively prime (i.e. $\gcd(m_i, m_j) = 1 \forall i \neq j$) integers ($\text{all } > 1$) & a_1, a_2, \dots, a_n are arbitrary integers, Then

$$x \equiv a_1 \pmod{m_1}$$

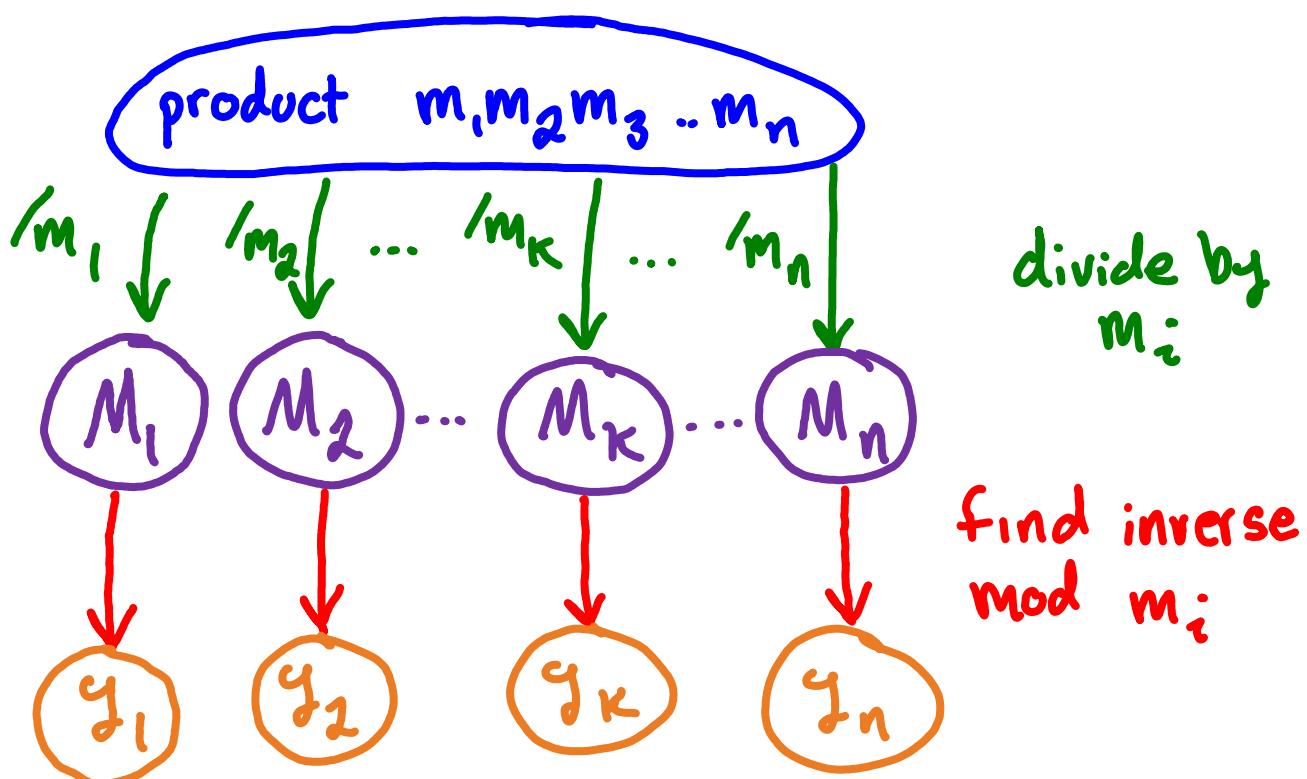
$$x \equiv a_2 \pmod{m_2}$$

⋮

$$x \equiv a_n \pmod{m_n}$$

Has a unique solution modulo $m = m_1 \cdot m_2 \cdot \dots \cdot m_n$
(i.e. \exists solution x with $0 \leq x < m$ & all other solutions
are congruent modulo m to this solution)

proof:



$$x = a_1 M_1 g_1 + a_2 M_2 g_2 + \dots + a_n M_n g_n$$

$$= \sum_{i=1}^n a_i M_i g_i$$

Check that x as defined above has the desired properties.

$$x \bmod m_1 =$$

$$\cancel{a_1 M_1 y_1 + a_2 M_2 y_2 + \dots + a_n M_n y_n} \bmod m_1$$

all terms except M_1 have a factor of m_1

thus $x \bmod m_1 = a_1 \cancel{M_1 y_1} \bmod m_1$

inverses modulo m_1

Therefore $x \bmod m_1 = a_1$

[i.e. $x \equiv a_1 \pmod{m_1}$] & similar

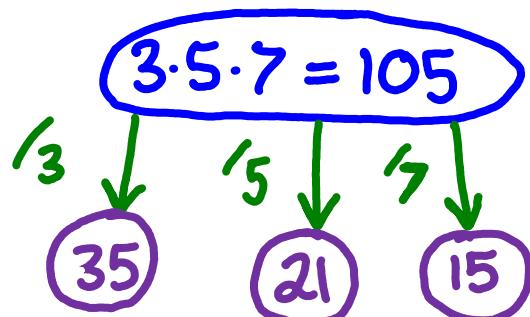
reasoning proves $x \equiv a_i \pmod{m_i} \forall i$. \square

Lets now revisit that example:

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

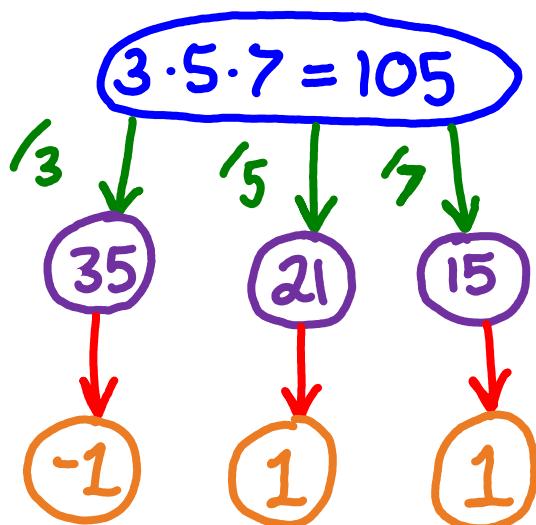
$$x \equiv 2 \pmod{7}$$



Next find inverses mod m_i

$$\begin{array}{l} \gcd(3, 35) = 1 \quad \& \quad 1 = 12 \cdot 3 - 1 \cdot 35 \\ \gcd(5, 21) = 1 \quad \& \quad 1 = 1 \cdot 21 - 4 \cdot 5 \\ \gcd(7, 15) = 1 \quad \& \quad 1 = 1 \cdot 15 - 2 \cdot 7 \end{array}$$

$$\begin{aligned} x &\equiv 2 \pmod{3} \\ x &\equiv 3 \pmod{5} \\ x &\equiv 2 \pmod{7} \end{aligned}$$



$$x := 2 \cdot 35 \cdot (-1) + 3 \cdot 21 \cdot 1 + 2 \cdot 15 \cdot 1$$

$$(\text{Simplify}) = -70 + 63 + 30 = \boxed{23}$$

Application: Computer Arithmetic w/ Large Integers

Computers generally run out of space to store numbers before running out of computing power. If we must carry out arithmetic w/ large numbers the Chinese Remainder Theorem (CRT) affords us the ability to exchange these numbers (which may be too big to store in the computer directly) for n-tuples of smaller numbers. The computer has an easier time computing with these n-tuples & the answer can be interpreted by translating back to a single large number by way of CRT.

(i) select moduli m_1, m_2, \dots, m_n

$\hookrightarrow m_i > 2 \quad \forall i \leq n$

$\hookrightarrow \gcd(m_i, m_j) = 1$ whenever $i \neq j$

$\hookrightarrow M := m_1 m_2 \cdots m_n$ is larger than all ~~**~~'s
we would like to do arithmetic with

(ii) Then $\forall x \leq M \exists!$ n -tuple of integers
representing x .

$$x \xrightarrow{\text{reduce}} (x \bmod m_1, x \bmod m_2, \dots, x \bmod m_n)$$

$\xleftarrow{\text{CRT}}$

(iii) Addition/Multiplication can be carried out
component-wise

$$\mathbb{Z}/M\mathbb{Z} \cong \mathbb{Z}/m_1\mathbb{Z} \times \mathbb{Z}/m_2\mathbb{Z} \times \cdots \times \mathbb{Z}/m_n\mathbb{Z}$$

provided the m_i are pairwise relatively prime

★ the above is an isomorphism of "rings"
(whatever that means)